

Boolean Functions for Cryptography and Coding Theory

Boolean functions are essential to systems for secure and reliable communication. This comprehensive survey of Boolean functions for cryptography and coding covers the whole domain and all important results, building on the authors influential articles with additional topics and recent results. A useful resource for researchers and graduate students, the book balances detailed discussions of properties and parameters with examples of various types of cryptographic attacks that motivate the consideration of these parameters. It provides all the necessary background on mathematics, cryptography, and coding and an overview of recent applications, such as side-channel attacks on smart cards and hardware, cloud computing through fully homomorphic encryption, and local pseudorandom generators. The result is a complete and accessible text on the state of the art in single- and multiple-output Boolean functions that illustrates the interaction among mathematics, computer science, and telecommunications.

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Preface

The present monograph is a merged, reorganized, significantly revised, and extensively completed version of two chapters, entitled "Boolean Functions for Cryptography and Error Correcting Codes" [236] and "Vectorial Boolean Functions for Cryptography" [237], which appeared in 2010 as parts of the book *Boolean Models and Methods in Mathematics, Computer Science, and Engineering* [394] (editors, Yves Crama and Peter Hammer). It is meant for researchers but is accessible to anyone who knows basics in linear algebra and general mathematics. All the other notions needed are introduced and studied (even finite fields are, in the Appendix).

Since these chapters were written in 2009, about 1,500 papers have been published that deal with this twofold topic (which is broad, as we see), and this version is updated with the main references and their main results (with corrections in the rare cases where they were needed). It also contains original results.

New notions on Boolean and vectorial functions and new ways of using them have also emerged. A chapter devoted to these recent and/or not enough studied directions of research has been included.

In the limit of a book, we tried to be as complete as possible. Of course, we could not go into details as much as do papers, but we made our best to ensure a good trade-off between completeness in scope and in depth. The choice of those papers that are referred to and of those results that are developed may seem subjective; it has been difficult, given the large number of papers. We tried, within the imposed length limit, to give the proof of a result each time it was short and simple enough, and when it provided a vision (we tried to avoid giving too technical proofs whose only – but of course important – value would have been to convince the reader that the result is true). We would have liked to avoid, when presenting arguments and observations, to refer to results (and concepts) to come later in the text, but the large number of results has made this necessary; otherwise, it would have been impossible to gather in a same place all the facts related to a same notion.

We have limited ourselves to Boolean and vectorial functions in characteristic 2, since these fit better with applications in coding and cryptography, and since dealing with p-ary and generalized functions would have reduced the description of the results on binary functions.



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Notation

```
|I|
                       size of a set I,
\lfloor u \rfloor
                       integer part (floor) of a real number u,
                       ceiling of u (the smallest integer larger than or equal to u),
\lceil u \rceil
\phi^{-1}(u)
                       preimage of u by a function \phi,
                       indicator (or characteristic) function of a set E: 1_E(x) = \begin{cases} 1 \text{ if } x \in E \\ 0 \text{ otherwise,} \end{cases}
1_E
\delta_a
                       the Dirac (or Kronecker) symbol at a (i.e. the indicator of \{a\}),
\mathbb{F}_2
                       the finite field with two elements 0, 1 (bits),
\mathbb{F}_2^n
                       the n-dimensional vector space over \mathbb{F}_2 (sometimes identified with \mathbb{F}_{2^n}),
\mathcal{L}_{n,m}
                       the vector space of linear (n, m)-functions,
0_n
                       zero vector in \mathbb{F}_2^n or in \mathbb{F}_q^n, n > 1 (in other groups, we just write 0),
1_n
                       vector (1, \ldots, 1) in \mathbb{F}_2^n,
                       addition in characteristic 0 (e.g., in \mathbb{R}), and in \mathbb{F}_2^n and \mathbb{F}_{2^n} for n > 1,
+
\sum_{i}
                       multiple sum of +,
                       addition in \mathbb{F}_2 (i.e., modulo 2); direct sum of two vector spaces,
\bigoplus_{i}
                       multiple sum of \oplus,
                       x + 1_n, where x \in \mathbb{F}_2^n,
a \cdot x
                       inner product in \mathbb{F}_2^n,
                       = a \cdot x, resp. x + a, where ":" is an inner product in \mathbb{F}_{2}^{n},
\ell_a(x), t_a(x)
\mathbb{F}_2^I
                       the vector space over \mathbb{F}_2 of all binary vectors whose indices range in I,
                       the finite (Galois) field of order 2^n, identified with \mathbb{F}_2^n as a vector space, = x + x^{2^m} + x^{2^{2m}} + \cdots + x^{2^{n-m}}, trace function from \mathbb{F}_{2^n} to \mathbb{F}_{2^m} (m \mid n), = tr_1^n(x) = \sum_{i=0}^{n-1} x^{2^i} the absolute trace function,
tr_m^n(x)
tr_n(x)
\mathbb{F}_{2^n}^*
                       \mathbb{F}_{2^n} \setminus \{0\}, where 0 denotes the zero element of \mathbb{F}_{2^n},
                       primitive element of \mathbb{F}_{2^n},
\alpha
                       convolutional product of two functions over \mathbb{F}_2^n (see page 60),
\otimes
f, g, h, \dots
                       Boolean functions,
\mathcal{BF}_n
                       the \mathbb{F}_2-vector space of all n-variable Boolean functions f: \mathbb{F}_2^n \to \mathbb{F}_2,
F, G, H, \dots
                       vectorial functions,
\mathcal{G}_F
                       graph of a vectorial function: \mathcal{G}_F = \{(x, F(x)); x \in \mathbb{F}_2^n\},
                       Hamming weight (of a vector, of a function),
w_H()
                       Hamming distance (between two vectors, two functions),
d_H(,)
d(C)
                       minimum (Hamming) distance of code C,
```



Notation xiii

```
supp()
                        the support (of a vector, of a function),
                         "x is covered by y" (i.e., supp(x) \subseteq supp(y)),
x \leq y
x \vee y
                         vector such that supp(x \lor y) = supp(x) \cup supp(y),
x \wedge y
                         vector such that supp(x \wedge y) = supp(x) \cap supp(y),
                         ith vector of the canonical basis of \mathbb{F}_2^n,
                        \prod_{i\in I} x_i, I\subseteq \{1,\ldots,n\}, \prod_{i=1}^n x_i^{u_i}, u\in \mathbb{F}_2^n,
                         binary Möbius transform (f^{\circ}: u \mapsto a_u, \text{ coef. of } x^u \text{ in the ANF of } f),
                         Fourier–Hadamard transform of a real-valued function \varphi over \mathbb{F}_2^n,
f_{\chi} W_f()
                         sign function of a Boolean function f, that is, x \mapsto (-1)^{f(x)},
                         Walsh transform of a Boolean function f(i.e., \widehat{f_{\gamma}}),
W_F(,)
                         Walsh transform of a vectorial function F,
supp(W_f)
                         support of W_f: \{u \in \mathbb{F}_2^n; W_f(u) \neq 0\},\
                         cardinality of the support of W_f,
N_{W_f}
                        \sum_{x \in \mathbb{F}_2^n} (-1)^{f(x)} (= W_f(0_n)),
\mathcal{F}(f)
nl()
                         nonlinearity of a Boolean or vectorial function,
                         r-th order nonlinearity of a Boolean function,
nl_r()
ln, log_2
                         natural (Neperian) logarithm, base 2 logarithm,
                         the algebraic degree of f (i.e., the degree of its ANF),
d_{alg}(f)
d_{num}(f)
                         the numerical degree of f (i.e., the degree of its NNF),
                         2-weight of integer j (see page 45),
w_2(j)
(n, m, t)-function
                         t-resilient (n, m)-function,
AI()
                         algebraic immunity of a function,
                        matrix of the system of equations \bigoplus_{I\subseteq\{1,\dots,n\}\atop III\subseteq J} a_I u^I = 0, u \in supp(f),
M_{f,d}
rk(M)
                         the rank of a matrix M,
FAC()
                         fast algebraic complexity of a function,
FAI()
                         fast algebraic immunity of a function,
                         derivatives in the direction a: x \mapsto f(x) \oplus f(x+a), F(x) + F(x+a),
D_a f, D_a F
                         the symmetric difference between two sets,
                         autocorrelation function \Delta_f(a) = \sum_{x \in \mathbb{F}_2^n} (-1)^{D_a f(x)},
\Delta_f(a)
                         absolute indicator of f: \Delta_f = \max_{a \in \mathbb{F}_2^n \setminus \{0_n\}} |\Delta_f(a)|,
\Delta_f
\mathcal{V}(f)
                         sum-of-squares indicator of f: \sum_{e \in \mathbb{F}_2^n} \mathcal{F}^2(D_e f),
                         linear kernel of a Boolean function f,
\mathcal{E}_f
RM(r,n)
                         Reed–Muller code of order r and length 2^n,
\rho(r,n)
                         covering radius of RM(r, n),
                         the symplectic form associated to a quadratic function f,
                         dual of a bent Boolean function (Definition 51, page 197),
\mathcal{M}
                         Maiorana-McFarland's class,
PS
                         partial spread class,
                         adjoint operator of a linear automorphism L,
Im(F)
                         the range (i.e., image set) F(\mathbb{F}_2^n) of an (n, m)-function,
                         the \mathbb{F}_2-vector space of annihilators of a Boolean function f,
An(f)
                         restriction of An(f) to those functions of algebraic degree at most d,
An_d(f)
B_{k,l}(f)
                         = \{g \in \mathcal{BF}_n; d_{alg}(g) \le k \text{ and } d_{alg}(fg) \le l\},\
```



xiv	Notation
f	defined by $f(x) = f(w_H(x))$, when f is symmetric,
$\sigma_i(x)$	elementary symmetric Boolean fct., of ANF: $\bigoplus_{I\subseteq\{1,,n\}/ I =i} x^I$,
$S_i(x)$	elementary symmetric pseudo-Boolean fct. NNF: $\sum_{I \subseteq \{1,,n\}/ I =i} x^I$,
δ_F	differential uniformity of an (n, m) -function F ,
Nb_F	imbalance of an (n, m) -function (see page 113),
NB_F	derivative imbalance of an (n, m) -function (see page 138),
X	a sharing of x (see page 436),
F	a threshold implementation of function F (see page 436),
$E_{n,k}$	$= \{x \in \mathbb{F}_2^n; \ w_H(x) = k\},$
$w_H(f)_k$	Hamming weight of the restriction of function f to $E_{n,k}$,