

## 1

## Introduction

In the first four sections we show how, starting with the usual description of free groups by means of reduced words, it is possible to arrive at a definition of the groups  $\mathcal{RF}(G)$  and their associated  $\mathbb{R}$ -trees  $\mathbf{X}_G$ , which are the objects of study in this book. The final section summarises the contents of the following chapters.

## 1.1 Finite words and free groups

In constructing free groups, one may start from the collection of all finite *words*

$$w = x_{i_1}^{e_1} x_{i_2}^{e_2} \cdots x_{i_n}^{e_n}$$

over an alphabet  $X \cup X^{-1}$ , where  $X$  is some given set,  $e_1, \dots, e_n \in \{1, -1\}$ , and

$$X^{-1} = \{x^{-1} : x \in X\}$$

is a set in one-to-one correspondence with  $X$  via the map  $x \mapsto x^{-1}$  such that  $X \cap X^{-1} = \emptyset$ . We extend this map to an involution of  $X \cup X^{-1}$  by setting  $(x^{-1})^{-1} = x$ . A word  $w$  can be thought of as a function

$$\{1, 2, \dots, n\} \rightarrow X \cup X^{-1},$$

for some integer  $n \geq 0$ , the unique word of length 0 being the *empty word*  $\varepsilon$ . A word  $w = x_{i_1}^{e_1} x_{i_2}^{e_2} \cdots x_{i_n}^{e_n}$  is called *reduced* if we have  $x_{i_j}^{e_j} \neq x_{i_{j+1}}^{-e_{j+1}}$  for all indices  $j$  with  $1 \leq j \leq n-1$ , that is, if  $w$  does not contain a subword of the form  $x_i^e x_i^{-e}$ . Clearly, the empty word  $\varepsilon$  itself is reduced.

## 1.2 Words over a discretely ordered abelian group $\Lambda$

One can generalise the above set-up by taking an arbitrary discretely ordered abelian group<sup>1</sup>  $\Lambda$ , and considering ‘infinite words’  $w : [1, \alpha] \rightarrow X \cup X^{-1}$  for  $\alpha \geq 0$ , where

$$[1, \alpha] = \{\beta \in \Lambda : 1 \leq \beta \leq \alpha\}$$

and where 1 denotes the least positive element of  $\Lambda$ , the case  $\alpha = 0$  corresponding to the empty word  $\varepsilon$ . This has indeed been done; see Myasnikov, Remeslenikov and Serbin [40]. In this setting the concept of reducedness still makes sense: a word  $w$  as above is *reduced*, if there does not exist  $\beta \in [1, \alpha - 1]$  such that  $w(\beta + 1) = w(\beta)^{-1}$ . Clearly, the empty word  $\varepsilon$  is reduced. Let  $R(\Lambda, X)$  be the set of all reduced words. We define the *inverse* of a word  $w$  on  $[1, \alpha]$  as the function  $w^{-1}$  given on the same domain  $[1, \alpha]$  by

$$w^{-1}(\beta) = w(\alpha - \beta + 1)^{-1}, \quad 1 \leq \beta \leq \alpha.$$

One can check immediately that if  $w$  is reduced then so is  $w^{-1}$ .

The concatenation of two words  $u, v$  on domains  $[1, \alpha]$  and  $[1, \beta]$ , respectively, is defined in a natural way as the word  $u \circ v$  with domain  $[1, \alpha + \beta]$  given by

$$(u \circ v)(\xi) = \left\{ \begin{array}{ll} u(\xi), & 1 \leq \xi \leq \alpha \\ v(\xi - \alpha), & \alpha + 1 \leq \xi \leq \alpha + \beta \end{array} \right\} \quad (\xi \in [1, \alpha + \beta]).$$

In this situation one can define a partial multiplication (reduced concatenation) on  $R(\Lambda, X)$  in a way that is analogous to multiplication in a free group. We first define, for  $u, v \in R(\Lambda, X)$ ,  $\text{com}(u, v)$  to be the largest common initial segment of  $u$  and  $v$ , more precisely,  $\text{com}(u, v) = u|_{[1, \gamma]}$  with  $\gamma \in \Lambda$  and  $\gamma \geq 0$  such that

$$u(\xi) = v(\xi), \quad \xi \in [1, \gamma],$$

and either  $\gamma = \min\{\alpha, \beta\}$  or  $u(\gamma + 1) \neq v(\gamma + 1)$ . The problem with this definition is, of course, that  $\text{com}(u, v)$  does not always exist, for which reason we shall only be able to define a partial multiplication on  $R(\Lambda, X)$ . Suppose that  $w := \text{com}(u^{-1}, v)$  is defined. Then we can write  $u^{-1} = w \circ u_1$ ,  $v = w \circ v_1$ , so that  $u = u_1^{-1} \circ w^{-1}$ , and we define the reduced product  $uv$  of the reduced words  $u$  and  $v$  by setting

$$uv = u_1^{-1} \circ v_1.$$

<sup>1</sup> By an ordered abelian group, we shall always mean a *totally ordered* abelian group.

1.2 Words over a discretely ordered abelian group  $\Lambda$  3

Since  $u$  and  $v$  are reduced, so is  $uv$ . In this way, we obtain a partial multiplication on  $R(\Lambda, X)$ , which one can show is associative if it is defined; that is, if  $uv$  and  $vw$  are defined, then  $(uv)w$  is defined if and only if  $u(vw)$  is defined, in which case  $(uv)w = u(vw)$ . (Unfortunately, none of the elegant constructions of a free group that circumvent the need for establishing associativity work directly in this situation.)

Note that the *empty word*  $\varepsilon$  (corresponding to  $\alpha = 0$ ) is a two-sided identity element, that is,

$$\varepsilon u = u = u\varepsilon, \quad u \in R(\Lambda, X).$$

Also, we have

$$uu^{-1} = \varepsilon = u^{-1}u, \quad u \in R(\Lambda, X).$$

Apart from the fact that reduced multiplication is only a partial operation, another marked difference from the free group case is that there can be words  $w$  with  $w \neq \varepsilon$  but  $w^2 = \varepsilon$ .

**Example 1.1** Let  $\Lambda = \mathbb{Z}^2$  with right lexicographic ordering, so that the least positive element is  $(1, 0)$ . Let  $\alpha = (0, 1)$  and fix  $x \in X$ . Define a word  $w$  on  $[(1, 0), \alpha]$  via

$$w(\beta) = \begin{cases} x, & \beta = (s, 0), s \geq 1, \\ x^{-1}, & \beta = (s, 1), s \leq 0. \end{cases}$$

Then  $w$  is reduced and non-trivial, and  $w^2 = \varepsilon$ .

There is also a notion of a *cyclically reduced* word: a word  $w \in R(\Lambda, X)$  is cyclically reduced if  $w(1) \neq w(\alpha)^{-1}$ . Let

$$\begin{aligned} CDR(\Lambda, X) &= \{w \in R(\Lambda, X) : w = u \circ v \circ u^{-1} \text{ for some cyclically reduced word } v\}. \end{aligned}$$

One can show that

$$CDR(\Lambda, X) = \{w \in R(\Lambda, X) : w^2 \text{ is defined and } w^2 \neq \varepsilon\} \cup \{\varepsilon\};$$

see Lemma 3.6 in [40].

We say that  $G \subseteq CDR(\Lambda, X)$  is a subgroup of  $CDR(\Lambda, X)$ , if  $u, v \in G$  implies that  $uv$  is defined and that  $uv \in G$ , if  $u \in G$  implies that  $u^{-1} \in G$ , and if  $\varepsilon \in G$ . If  $G$  is a subgroup of  $CDR(\Lambda, X)$ , one can show that the function  $L : G \rightarrow \Lambda$  given by  $L(w) = \alpha$ , where the domain of  $w$  is  $[1, \alpha]$ , and  $L(\varepsilon) = 0$ , is a Lyndon length function on  $G$  and gives rise to an action of  $G$  on a  $\Lambda$ -tree that is

free and without inversions. (These terms are explained in Appendix A.) This generalises the fact that a free group, and so any subgroup, acts freely on its Cayley graph with respect to a basis; this graph is a tree. In fact, one can prove the following.

**Theorem 1.2** *Let  $\Lambda$  be a discretely ordered abelian group. A group  $G$  acts freely and without inversions on a  $\Lambda$ -tree if and only if  $G$  is a subgroup of  $CDR(\Lambda, X)$  for some set  $X$ .*

This is shown in [11]; the backward implication also appears in [40].

### 1.3 The case where $\Lambda$ is densely ordered

At this stage, the question arises: can something analogous be done if instead we start from a *densely ordered* abelian group  $\Lambda$ ? The first problem is that there is no longer a least positive element, so we replace a domain  $[1, \alpha]$  with an interval  $[0, \alpha]$  where  $\alpha \geq 0$ . A more serious problem, however, is that concatenation can no longer be defined as above. Our solution is to replace the set  $X \cup X^{-1}$  by a (discrete) group  $G$ . Let

$$\mathcal{F}(\Lambda, G) := \bigcup_{\substack{\alpha \in \Lambda \\ \alpha \geq 0}} G^{[0, \alpha]} = \{f : [0, \alpha] \rightarrow G : \alpha \in \Lambda, \alpha \geq 0\}$$

be the set of all functions with values in  $G$  defined on an interval of  $\Lambda$  of the form  $[0, \alpha]$  for some  $\alpha \geq 0$ . Concatenation is then replaced by an operation denoted  $*$ , the *star product*, defined as follows: if  $f, g \in \mathcal{F}(\Lambda, G)$  are functions with domains  $[0, \alpha]$  and  $[0, \beta]$ , respectively, then  $f * g$  is the function given on the interval  $[0, \alpha + \beta]$  of  $\Lambda$  via

$$(f * g)(\xi) = \left. \begin{cases} f(\xi), & 0 \leq \xi < \alpha \\ f(\alpha)g(0), & \xi = \alpha \\ g(\xi - \alpha), & \alpha < \xi \leq \alpha + \beta \end{cases} \right\} (\xi \in [0, \alpha + \beta]).$$

The function  $\mathbf{1}_G$  defined on the interval  $[0, 0] = \{0\}$  by  $\mathbf{1}_G(0) = 1_G$  (where  $1_G$  is the identity element of  $G$ ) is a two-sided identity element with respect to the star operation; that is, we have

$$f * \mathbf{1}_G = f = \mathbf{1}_G * f, \quad f \in \mathcal{F}(\Lambda, G).$$

1.4 The case where  $\Lambda = \mathbb{R}$  5

We also have a notion of the *formal inverse*  $f^{-1}$  of a function  $f \in \mathcal{F}(\Lambda, G)$ : if  $f$  is defined on the domain  $[0, \alpha]$  then  $f^{-1}$  is the function given on the same interval  $[0, \alpha]$  by

$$f^{-1}(\xi) = (f(\alpha - \xi))^{-1}, \quad 0 \leq \xi \leq \alpha.$$

In this setting there is also a notion of a reduced function, which necessarily needs to be somewhat more elaborate. A function  $f \in \mathcal{F}(\Lambda, G)$  defined on the interval  $[0, \alpha]$  of  $\Lambda$  is called *reduced* if, for each point  $\xi_0 \in (0, \alpha)$  with  $f(\xi_0) = 1_G$  and every element  $\varepsilon \in \Lambda$  with  $0 < \varepsilon \leq \min\{\alpha - \xi_0, \xi_0\}$ , there exists some  $\delta \in \Lambda$  such that  $0 < \delta \leq \varepsilon$  and such that  $f(\xi_0 + \delta) \neq (f(\xi_0 - \delta))^{-1}$ . The set of all reduced functions in  $\mathcal{F}(\Lambda, G)$  is denoted by  $\mathcal{RF}(\Lambda, G)$ . Given a function  $f : [0, \alpha] \rightarrow G$  in  $\mathcal{F}(\Lambda, G)$ , let us call an  $\varepsilon$ -neighbourhood

$$[\xi_0 - \varepsilon, \xi_0 + \varepsilon] \subseteq [0, \alpha]$$

of a point  $\xi_0 \in (0, \alpha)$ , with  $f(\xi_0) = 1_G$ , a *cancelling neighbourhood around  $\xi_0$*  if  $f(\xi_0 - \delta) = (f(\xi_0 + \delta))^{-1}$  for all  $0 < \delta \leq \varepsilon$ . Then we can say that a function  $f \in \mathcal{F}(\Lambda, G)$  as above is reduced if and only if there does not exist a cancelling neighbourhood around any interior point of the domain  $[0, \alpha]$  of  $f$  satisfying  $f(\xi_0) = 1_G$ .

For  $u, v \in \mathcal{F}(\Lambda, G)$ , an analogue of  $\text{com}(u, v)$  can be defined and, if the element  $\text{com}(u^{-1}, v) =: w$  exists, so that  $u^{-1} = w * u_1$  and  $v = w * v_1$ , we may define the reduced product  $uv$  of  $u$  and  $v$  by  $uv = u_1^{-1} * v_1$ . This gives a partial multiplication on  $\mathcal{F}(\Lambda, G)$  that is associative when defined. It can also be shown that the product of two reduced functions, when it exists, is again reduced.

1.4 The case where  $\Lambda = \mathbb{R}$

In this book, we shall confine our attention to the case where  $\Lambda = \mathbb{R}$ , taking the view that this is already quite difficult to deal with (in particular, the proof of the associativity of reduced multiplication is non-trivial). In this case (reduced) multiplication is always defined, and we obtain a group denoted by  $\mathcal{RF}(G)$ , with the formal inverse of a reduced function  $f$  acting as a two-sided inverse of  $f$  and with  $\mathbf{1}_G$  as the neutral element.

There is a construction of an  $\mathbb{R}$ -tree  $\mathbf{X}_G$  on which  $\mathcal{RF}(G)$  acts with point stabilisers isomorphic to  $G$ . More precisely, by definition, each element  $f$  of  $\mathcal{RF}(G)$  has a real number  $L(f)$  assigned to it, namely the length  $\alpha$  of its

domain  $[0, \alpha]$ ; it is not hard to see that the function  $L : \mathcal{RF}(G) \rightarrow \mathbb{R}$  defined in this way is a Lyndon length function. It follows that  $\mathcal{RF}(G)$  has a canonical action by isometries on an  $\mathbb{R}$ -tree

$$\mathbf{X}_G = (X_G, d)$$

with a distinguished base-point  $x_0$  such that  $L_{x_0} = L$ , where  $L_{x_0}$  is the displacement function

$$L_{x_0}(f) = d(x_0, fx_0), \quad f \in \mathcal{RF}(G)$$

associated with this action and such that

$$G_0 := \text{stab}_{\mathcal{RF}(G)}(x_0) = \{f \in \mathcal{RF}(G) : L(f) = 0\} \cong G.$$

It turns out that  $\mathbf{X}_G$  is always metrically complete and that the action of  $\mathcal{RF}(G)$  on  $\mathbf{X}_G$  is transitive.

As always in such situations, the action of  $\mathcal{RF}(G)$  on  $\mathbf{X}_G$  leads to a classification of the elements of  $\mathcal{RF}(G)$  according to whether they are elliptic (that is, have a fixed point) or hyperbolic (that is, act as a fixed-point free isometry). Hyperbolic elements have some local geometry associated with them, leading, in particular, to another type of length function on  $\mathcal{RF}(G)$ : if  $f \in \mathcal{RF}(G)$  is hyperbolic then there exists an isometric copy  $A_f \subseteq \mathbf{X}_G$  of the real line (the so-called *axis* of  $f$ ) such that  $f$  acts on  $A_f$  as a non-trivial translation; in particular, hyperbolic elements have infinite order. The translation length of a hyperbolic element  $f$  along its axis  $A_f$  is called the *hyperbolic length* of  $f$ , denoted  $\ell(f)$ , and  $\ell$  is extended to the whole of  $\mathcal{RF}(G)$  by setting  $\ell(f) = 0$  for an elliptic function  $f$ .

With a view to investigating further the action of  $\mathcal{RF}(G)$ , we shall introduce and study an analogue of cyclic reduction in free groups; this allows us, among other things, to characterize hyperbolic elements in a purely algebraic way and to compute hyperbolic length in terms of the length function  $L$ .

It follows from the transitivity of the action that the set of elliptic elements in  $\mathcal{RF}(G)$  coincides with the union of all conjugates of  $G_0$ ; in particular,  $\mathcal{RF}(G)$  has torsion if and only if the group  $G$  has. We will establish a stronger result to the effect that a subgroup of  $\mathcal{RF}(G)$  is bounded (with respect to the length function  $L$ ) if and only if it is conjugate to a subgroup of  $G_0$ . This result shows in particular that every finite subgroup of  $\mathcal{RF}(G)$  is conjugate to a subgroup of  $G_0$ , a result reminiscent of the bounded subgroup theorem for free products with amalgamation; see, for instance, Theorem 8 of Chapter I in Serre [45]. It also follows that the trivial group  $\{1_G\}$  is the only bounded subnormal subgroup of  $\mathcal{RF}(G)$ .

**1.5 Contents of the book**

We now turn to a discussion of individual chapters.

*Chapters 2 and 3.* Here we give the basic definitions, introduce the groups  $\mathcal{RF}(G)$ , and develop some cancellation theory needed for (among other things) a proof of the associativity of reduced multiplication. We then study the geometry associated with  $\mathcal{RF}(G)$  via its action on the  $\mathbb{R}$ -tree  $\mathbf{X}_G$  and the classification of the group elements effected by this action, covering the ground indicated in Section 1.4 above and more.

*Chapter 4.* This chapter reflects a rather exciting new development, reporting on the authors' recent discovery of two basic *embedding theorems*. We show that a group  $G$  acting freely and without inversions on a  $\Lambda$ -tree  $\mathbf{X}$  (for an arbitrary ordered abelian group  $\Lambda$ ) can be embedded into a group  $\hat{G}$ , acting freely, without inversions, and *transitively* on the  $\Lambda$ -tree  $\hat{\mathbf{X}}$ , which isometrically and  $G$ -equivariantly embeds  $\mathbf{X}$ . This result is of considerable independent interest, in particular shedding new light on the class of infinitely generated  $\mathbb{R}$ -free groups.

We then proceed to discuss a second, more specialised, embedding theorem concerning free and transitive  $\mathbb{R}$ -tree actions: we show that a group  $G$  acting freely and transitively on an  $\mathbb{R}$ -tree  $\mathbf{X}$  can be embedded into  $\mathcal{RF}(H)$  for some suitable group  $H$  such that  $\mathbf{X}$  embeds isometrically and  $G$ -equivariantly into  $\mathbf{X}_H$ , the  $\mathbb{R}$ -tree canonically associated with  $\mathcal{RF}(H)$ .

Combining these two results we infer that  $\mathcal{RF}$ -groups and their associated  $\mathbb{R}$ -trees are in fact *universal* (with respect to inclusion) for free  $\mathbb{R}$ -tree actions.

*Chapter 5.* Very little is known at present concerning homomorphisms involving  $\mathcal{RF}$ -groups. In this chapter a certain homomorphism

$$e_g : \mathcal{RF}(G) \rightarrow \mathbb{R}$$

is defined for each element  $g \in G$  by means of Lebesgue measure theory. The construction of these maps  $e_g$  is analogous to and inspired by the exponent sum maps of a free group relative to a basis element. By construction, the elliptic elements of  $\mathcal{RF}(G)$  are contained in the kernel of  $e_g$  for every  $g$ , and if  $g \in G$  is not an involution then the corresponding map  $e_g$  is surjective; this shows in particular that if  $G$  is not an elementary abelian 2-group then  $\mathcal{RF}(G)$  is not generated by its elliptic elements. For  $G$  an elementary abelian 2-group, the question remains open at this stage since all exponent sums of  $G$  are trivial.

The problem is taken up and resolved in Chapter 9 as part of the theory of test functions (see below).

*Chapter 6.* In this chapter we explore various aspects of functoriality of the  $\mathcal{RF}$ -construction. The most striking result obtained here is that if two groups  $G$  and  $H$  have the same (cardinal) number of involutions and the same number of non-involutions then we have

$$\mathcal{RF}(G)/E(G) \cong \mathcal{RF}(H)/E(H),$$

where  $E(G)$  is the subgroup of  $\mathcal{RF}(G)$  generated by the elliptic elements, that is, the normal closure of  $G_0$ . With slight imprecision the last result may be rephrased as follows.

*The isomorphism type of the group  $\mathcal{RF}(G)/E(G)$  depends only on the two cardinal numbers  $|\text{Inv}(G)|$  and  $|G - \text{Inv}(G)|$ .*

The proof of this surprising and rather deep lying *rigidity result* is long and somewhat technical. However, the techniques developed in this chapter also allow us to obtain at least a partial result concerning the automorphism group of the quotient group  $\mathcal{RF}(G)/E(G)$ ; see Proposition 6.7 and Corollary 6.8.

*Chapter 7.* A guiding principle when investigating  $\mathcal{RF}$ -groups appears to be the following.

*Hyperbolic elements of a non-trivial  $\mathcal{RF}$ -group behave analogously to the non-trivial elements of a (large) free group.*

This principle manifests itself for instance in the *conjugacy theorem for hyperbolic elements* established in Chapter 7, which (except for its proof) is an exact continuous analogue of the corresponding result for free groups. We also show there that the centraliser of a hyperbolic element  $f \in \mathcal{RF}(G)$  has index at most 2 in the normaliser of the infinite cyclic group  $\langle f \rangle$  in  $\mathcal{RF}(G)$ .

*Chapter 8.* It is easy to see that, for a non-trivial element  $g \in G_0$ , we have

$$C_{\mathcal{RF}(G)}(g) = C_{G_0}(g).$$

Consequently the centralisers of elliptic elements in  $\mathcal{RF}(G)$  are determined, up to isomorphism, by the isomorphism types of the centralisers in the group  $G$  itself; hence, in general (that is, without restricting the structure of  $G$ ), nothing more can be said here.

The situation is very different, and much more interesting, for hyperbolic elements, and the present chapter provides a penetrating study of their centralisers. We establish a criterion characterising those hyperbolic elements whose



centraliser is cyclic and obtain considerable insight into the centraliser structure in the general case; in particular, we show that centralisers of hyperbolic elements are abelian and relatively ‘small’, in that they always embed into the additive reals. As suggested by Remeslennikov, the centraliser  $C_{\mathcal{RF}(G)}(f)$  for hyperbolic  $f$  is controlled by (a subset of) the periods of the function  $f$ ; the reader is referred to Chapter 8 for details.

As an application of the main result of that chapter (Theorem 8.16), we show that  $\mathcal{RF}$ -groups enjoy an analogue of the *centraliser partition property* of free groups: the binary relation  $\leftrightarrow$  given by

$$f \leftrightarrow g : \iff f \text{ and } g \text{ commute}$$

is an equivalence relation on the set

$$\mathcal{RF}(G) = \bigcup_{t \in \mathcal{RF}(G)} tG_0t^{-1}$$

of hyperbolic elements of  $\mathcal{RF}(G)$ ; see Proposition 8.23 and Corollary 8.24. This result provides a further illustration of the philosophy concerning hyperbolic elements expressed above. As another application of Theorem 8.16, we show that  $\mathcal{RF}$ -groups do not contain non-trivial soluble normal subgroups. A completely different approach to this last result, using the theory of test functions, is given in Chapter 10.

*Chapters 9 and 10.* These two chapters provide an introduction to the theory of test functions and its applications, as developed originally in Müller [36] and Müller and Schlage-Puchta [38]. Roughly speaking, a test function is a mapping  $f : [0, \alpha] \rightarrow G$  of positive length  $L(f) = \alpha$ , such that  $f$  does not look locally like its own inverse. More precisely, we require that there do not exist  $\varepsilon > 0$  and points  $\xi_1, \xi_2 \in (0, \alpha)$  such that

$$f(\xi_1 + \eta) = f^{-1}(\xi_2 + \eta), \quad |\eta| < \varepsilon.$$

Test functions do in fact always exist; for instance, the function  $f_0$  of length 1 given by

$$f_0(\xi) = \begin{cases} x, & \xi^2 \in \mathbb{Q} \\ 1_G, & \xi^2 \notin \mathbb{Q} \end{cases} \quad (0 \leq \xi \leq 1),$$

where  $x$  is any non-trivial element of  $G$ , can be shown to be a test function; see Section 9.3. Test functions are automatically (cyclically) reduced and give rise to a further class of homomorphisms  $\mathcal{RF}(G) \rightarrow \mathbb{R}$ . Roughly speaking, given a test function  $f \in \mathcal{RF}(G)$  of length  $\alpha > 0$ , the idea is to compare (‘test’)

functions  $g \in \mathcal{F}(G)$  locally against  $f$  and  $f^{-1}$ , in this way obtaining two sets  $\mathcal{M}_f^+(G), \mathcal{M}_f^-(G) \subseteq (0, L(g))$ . To be more explicit, we set

$$\mathcal{M}_f^\pm(g) := \left\{ \xi \in (0, L(g)) : \exists \varepsilon > 0, \exists \xi' \in (0, \alpha) \text{ such that} \right. \\ \left. g(\xi + \eta) = f^\pm(\xi' + \eta) \text{ for all } |\eta| < \varepsilon \right\},$$

observing that  $\mathcal{M}_f^+(g)$  and  $\mathcal{M}_f^-(g)$  are open sets and thus Lebesgue measurable, and define a function  $\lambda_f : \mathcal{RF}(G) \rightarrow \mathbb{R}$  by

$$\lambda_f(g) = \mu(\mathcal{M}_f^+(g)) - \mu(\mathcal{M}_f^-(g)),$$

where  $\mu$  denotes Lebesgue measure. We show that  $\lambda_f$  is a surjective homomorphism whose kernel contains  $E(G)$ , in this way demonstrating in particular that  $\mathcal{RF}(G)$  is never generated by its elliptic elements; see Theorem 9.8 and Corollary 9.9.

A second important idea introduced in Chapter 9 is that of *local compatibility* and *incompatibility*. Roughly speaking, given functions  $f : [0, \alpha] \rightarrow G$  and  $g : [0, \beta] \rightarrow G$ , we say that  $f$  and  $g$  are locally compatible if  $f$  looks locally like  $g$  or  $g^{-1}$ . To be more precise,  $f$  and  $g$  as above are termed locally compatible if there exist  $\varepsilon > 0$  and points  $\xi \in (0, \alpha), \zeta \in (0, \beta)$  such that either

$$f(\xi + \eta) = g(\zeta + \eta), \quad |\eta| < \varepsilon,$$

or

$$f(\xi + \eta) = g^{-1}(\zeta + \eta), \quad |\eta| < \varepsilon.$$

If  $f$  and  $g$  both have positive length but are not locally compatible then they are called locally incompatible. Locally incompatible functions have no cancellation against each other, and if  $f, g$  are locally incompatible then so are  $f^{-1}$  and  $g$  as well as  $f^{-1}$  and  $g^{-1}$ .

We call a subgroup  $\mathcal{H} \leq \mathcal{RF}(G)$  *hyperbolic* if the set  $\mathcal{H} - \{\mathbf{1}_G\}$  consists entirely of hyperbolic elements. As a further application of test function theory (as developed so far), in Section 9.6 we show among other things that the family of centralisers  $\{C_{\mathcal{RF}(G)}(f_\sigma)\}_{\sigma \in S}$  corresponding to a family  $\{f_\sigma\}_{\sigma \in S}$  of pairwise locally incompatible test functions generates a hyperbolic subgroup of  $\mathcal{RF}(G)$  isomorphic to the free product  $\ast_{\sigma \in S} C_{\mathcal{RF}(G)}(f_\sigma)$ ; see Corollary 9.24.

The most striking applications of test function theory to date, however, stem from a rather deep result (Theorem 10.1) asserting the existence of large families of pairwise locally incompatible test functions with prescribed centraliser: *given a non-trivial group  $G$  and a proper subgroup  $\Lambda \leq (\mathbb{R}, +)$ , there exists a family  $\mathfrak{F}$  of pairwise locally incompatible test functions in  $\mathcal{RF}(G)$  such*